MIXED INTEGER LINEAR PROGRAMMING PROBLEMS WITH FUZZY VARIABLES

S. C. Sharma* and Abha Bansal**

* Department of Mathematics, University of Rajasthan, Jaipur-302055, India.
E-mail: sureshchand26@gmail.com

** Lakshya Institute, Opp. Head Post Office, New Colony, Dungurpur -314001, Rajasthan, India.
E-mail: neelamchikusinghal2002@gmail.com

(Received on: 25-02-11; Accepted on: 05-03-11)

------------------------------------------------------------------------------------------------------------------------------------------------

ABSTRACT

In this paper, a new decomposition method is used for solving mixed integer linear programming problems with fuzzy variables by using classical mixed integer linear programming has been proposed. In the decomposition method, ranking functions are not used. This method provides the best solution to a variety of mixed integer linear programming problems with fuzzy variables in a simple and effective manner. The decomposition method is illustrated with the help of a numerical example.

Key Words: Linear Programming, Integer Programming, Fuzzy Variables.

------------------------------------------------------------------------------------------------------------------------------------------------

1. INTRODUCTION:

Mixed integer linear programming is linear programming which require some or all of the variables to be integers. In real world situation the available information in the system under consideration are not exact, therefore fuzzy linear programming (FLP) was introduced and studied by many researchers, see, e.g., [3, 4, 6, 8, 11]. Fuzzy set theory has been applied to many disciplines such as control theory and management sciences, mathematical modeling and industrial applications. The concept of FLP on a general level was first proposed by Tanaka et al. [10].

FLP problems have an essential role in fuzzy modeling, which can formulate uncertainty in actual environment. Afterwards, many authors have considered various types of the FLP problems and proposed several approaches for solving these problems. In particular, the most convenient methods are based on the concept of comparison of fuzzy numbers with help of ranking functions, see, e.g., [3, 8]. Usually in such methods authors define a crisp model which is equivalent to the FLP problem and then use optimal solution of the model as the optimal solution of the FLP problem.

Mahdavi-Amiri and Nasseri [7] introduced a dual simplex algorithm for solving linear programming problem with fuzzy variables and its dual by using a general linear ranking function and linear programming directly. Allahviranloo et al. [1] have proposed a new method based on fuzzy number ranking function for a FILP problem via crisp integer linear programming (ILP) problems. Nasseri [9] has proposed a new method for solving the FLP problems in which he has used the fuzzy ranking method for converting the fuzzy objective function into crisp objective function.

In this paper, we have proposed a new method namely, decomposition method for solving a MILP problem with fuzzy variables by using the classical MILP. The significance of this paper is providing a new method for solving MILP problems with fuzzy variables without using any ranking functions. This method can serve managers by providing the best solution to a variety of mixed integer linear programming problems with fuzzy variables in a simple and effective manner. The decomposition method is illustrated with the help of numerical example.

2. PRELIMINARIES:

We need the following definitions of the basic arithmetic operators on fuzzy triangular numbers based on the function principle which can be found in [9].

Definition 2.1. A fuzzy number $\tilde{A}$ is a triangular fuzzy number denoted by $(a_1, a_2, a_3)$ where $a_1, a_2$ and $a_3$ are real numbers and its membership function is given below.

$$
\mu_\tilde{A}(x) = \begin{cases} 
\frac{(x - a_1)}{(a_2 - a_1)} & \text{for } a_1 \leq x \leq a_2 \\
\frac{(a_3 - x)}{(a_3 - a_2)} & \text{for } a_2 \leq x \leq a_3 \\
0 & \text{otherwise}
\end{cases}
$$
Definition: 2.2. Let \((a_1, a_2, a_3)\) be two triangular fuzzy numbers. Then

(a) \((a_1, a_2, a_3) + (b_1, b_2, b_3) = (a_1 + b_1, a_2 + b_2, a_3 + b_3)\)
(b) \((a_1, a_2, a_3) - (b_1, b_2, b_3) = (a_1 - b_1, a_2 - b_2, a_3 - b_3)\)
(c) \(k(a_1, a_2, a_3) = (ka_1, ka_2, ka_3)\) for \(k \geq 0\)
(d) \(k(a_1, a_2, a_3) = (ka_3, ka_2, ka_1)\) for \(k < 0\)

Definition: 2.3. Let \(\tilde{A} = (a_1, a_2, a_3)\) be in \(F(R)\). Then

(i) \(\tilde{A}\) is said to be positive if \(a_i \geq 0\), for all \(i = 1\) to 3;
(ii) \(\tilde{A}\) is said to be integer if \(a_i\) for all \(i = 1\) to 3 are integers and
(iii) \(\tilde{A}\) is said to be symmetric if \(a_2 - a_1 = a_3 - a_2\).

Definition: 2.4. A fuzzy vector \(\tilde{b} = (\tilde{b}_j)_{m \times n}\) is called nonnegative and denoted by \(\tilde{b} \geq 0\), if each element of \(\tilde{b}\) is a nonnegative real fuzzy number, that is, \(\tilde{b}_j \geq 0\), \(i = 1, 2, \ldots, m\).

Consider the following \(m \times n\) fuzzy linear system with nonnegative real fuzzy numbers:

\[A \tilde{x} \leq \tilde{b}\]

Where \(A = (a_{ij})_{m \times n}\) is a nonnegative crisp matrix and \(\tilde{x} = (\tilde{x}_j), \tilde{b} = (\tilde{b}_j)\) are nonnegative fuzzy vectors and \(\tilde{x}_j, \tilde{b}_j \in F(R)\) for all \(1 \leq j \leq n\) and \(1 \leq i \leq m\).

3. FUZZY MIXED INTEGER LINEAR PROGRAMMING:

Consider the following integer linear programming problem with fuzzy variables

\[
\max \tilde{z} = \tilde{c}\tilde{x} \\
\text{subject to } A\tilde{x} \leq \tilde{b}, \quad \tilde{x} \geq 0 \quad \text{and are integers}
\]

Where the coefficient matrix \(A = (a_{ij})_{m \times n}\) is a nonnegative real crisp matrix.

The cost vector \(c = (c_1, \ldots, c_n)\) is nonnegative crisp vector and \(\tilde{x} = (\tilde{x}_j)_{n \times 1}\) and \(\tilde{b} = (\tilde{b}_j)_{m \times 1}\) are nonnegative real fuzzy vectors such that \(\tilde{x}_j, \tilde{b}_j \in F(R)\) for all \(1 \leq j \leq n\) and \(1 \leq i \leq m\). If some or all variables in \(\tilde{x} = (\tilde{x}_j)_{n \times 1}\) are restricted to be integer then the problem is called fuzzy mixed integer linear programming problem.

4. THE ALGORITHM:

Step 1: Consider the mixed integer linear programming problem with fuzzy variables

\[
\max \tilde{z} \\
\text{subject to } \tilde{x}_1 \geq 0 \text{ and integer, } \tilde{x}_2 \geq 0.
\]

Step 2: Let \(\tilde{z} = (z_1, z_2, z_3), \tilde{x}_1 = (y_1, x_1, t_1)\) and \(\tilde{x}_2 = (y_2, x_2, t_2)\).

Step 3: Then the above problem is converted into three crisp mixed integer linear programming problem.

\[
(P_2) \max_{x_1, x_2} z_2 \\
\text{subject to } x_1 \geq 0 \text{ and integer } x_2 \geq 0.
\]
Step 4: Introduce slack or surplus variables to convert the inequalities into equations and obtain the optimum solution of given problem by the dual simplex method.

Step 5: Let \((u_1, u_2)\) be its optimal solution.

Step 6: Test the integrality condition of the optimum.
   
i. If \(u_1 \geq 0\) and are integer, \(u_2 \geq 0\), the current solution is an optimal one then goto Step 16.
   
ii. If \(u_1 \geq 0\) but not integer then go to the next step.

Step 7: Express the number in two parts as the sum of the integer and a non-negative fraction.

Step 8: Apply the tree of Branch and Bound method to get the integer solution.

Step 9: Let \((U_1, U_2)\) be its optimal solution.

Step 10: Write the second mixed integer linear programming problem.

\[(P_2) \quad \max_{y_1, y_2} z_1\]

subject to \(y_1 \leq U_1\)
\(y_2 \leq U_2\)
\(y_1 \geq 0\) and integer \(y_2 \geq 0\).

Step 11: Go to step 4 to step 8 and repeat the procedure until we get the optimum solution which satisfy the integrality condition.

Step 12: Let \((V_1, V_2)\) be its optimal solution.

Step 13: Write the third mixed integer linear programming problem.

\[(P_3) \quad \max_{t_1, t_2} z_3\]

subject to \(t_1 \leq V_1\)
\(t_2 \leq V_2\)
\(t_1 \geq 0\) and integer \(t_2 \geq 0\).

Step 14: Go to step 4 to step 8 and repeat the procedure until we get the optimum solution which satisfy the integrality condition.

Step 15: Let \((W_1, W_2)\) be its optimal solution.

Step 16: We get the final required optimal solution is \(\tilde{x}_1 = (y_1, x_1, t_1) = (U_1, V_1, W_1)\)
\(\tilde{x}_2 = (y_2, x_2, t_2) = (U_2, V_2, W_2)\), \(\tilde{z} = (z_1, z_2, z_3)\)

5. NUMERICAL EXAMPLE:

The proposed method is illustrated by the following example.

Consider the following integer linear programming problem with

\[
\max \tilde{z} = 5\tilde{x}_1 \oplus 6\tilde{x}_2\]

subject to \(10\tilde{x}_1 \oplus 3\tilde{x}_2 = (48, 52, 48)\)
\(2\tilde{x}_1 \oplus 3\tilde{x}_2 = (12, 18, 12)\)
\(\tilde{x}_1 \geq 0\) and integer
\[ x_2 \geq 0 \]

**Solution:**
Let \( \tilde{z} = (\tilde{z}_1, \tilde{z}_2, \tilde{z}_3) \), \( \tilde{x}_1 = (y_1, x_1, t_1) \) and \( \tilde{x}_2 = (y_2, x_2, t_2) \)

Now, the problem \((P_2)\) is given below

\[
\begin{align*}
\text{(P}_2\text{)} & \quad \text{max } z_2 = 5x_1 + 6x_2 \\
\text{subject to } & \quad 10x_1 + 3x_2 \leq 52 \\
& \quad 2x_1 + 3x_2 \leq 18 \\
& \quad x_1 \geq 0 \text{ and integer} \\
& \quad x_2 \geq 0
\end{align*}
\]

Now, using dual simplex method

**Table-1**

<table>
<thead>
<tr>
<th>Basic variables</th>
<th>( C_j )</th>
<th>5</th>
<th>6</th>
<th>0</th>
<th>0</th>
<th>Min ( X_B / X_k )</th>
</tr>
</thead>
<tbody>
<tr>
<td>( C_B )</td>
<td>( X_B )</td>
<td>( x_1 )</td>
<td>( x_2 )</td>
<td>( x_3 )</td>
<td>( x_4 )</td>
<td></td>
</tr>
<tr>
<td>( x_3 )</td>
<td>0</td>
<td>52</td>
<td>10</td>
<td>3</td>
<td>1</td>
<td>0</td>
</tr>
<tr>
<td>( x_4 )</td>
<td>0</td>
<td>18</td>
<td>2</td>
<td>3</td>
<td>0</td>
<td>1</td>
</tr>
<tr>
<td>( Z_j - C_j )</td>
<td>-5</td>
<td>-6</td>
<td>0</td>
<td>0</td>
<td></td>
<td></td>
</tr>
</tbody>
</table>

**Table-2**

<table>
<thead>
<tr>
<th>Basic variables</th>
<th>( C_j )</th>
<th>5</th>
<th>6</th>
<th>0</th>
<th>0</th>
<th>Min ( X_B / X_k )</th>
</tr>
</thead>
<tbody>
<tr>
<td>( C_B )</td>
<td>( X_B )</td>
<td>( x_1 )</td>
<td>( x_2 )</td>
<td>( x_3 )</td>
<td>( x_4 )</td>
<td></td>
</tr>
<tr>
<td>( x_3 )</td>
<td>0</td>
<td>34</td>
<td>8</td>
<td>0</td>
<td>1</td>
<td>-1</td>
</tr>
<tr>
<td>( x_2 )</td>
<td>6</td>
<td>6</td>
<td>2/3</td>
<td>1</td>
<td>0</td>
<td>1/3</td>
</tr>
<tr>
<td>( Z_j - C_j )</td>
<td>-1</td>
<td>0</td>
<td>0</td>
<td>2</td>
<td></td>
<td></td>
</tr>
</tbody>
</table>

Since all \( Z_j - C_j \geq 0 \), thus Table-3 is an optimal one.

**Table-3**

<table>
<thead>
<tr>
<th>Basic variables</th>
<th>( C_j )</th>
<th>5</th>
<th>6</th>
<th>0</th>
<th>0</th>
<th>Min ( X_B / X_k )</th>
</tr>
</thead>
<tbody>
<tr>
<td>( C_B )</td>
<td>( X_B )</td>
<td>( x_1 )</td>
<td>( x_2 )</td>
<td>( x_3 )</td>
<td>( x_4 )</td>
<td></td>
</tr>
<tr>
<td>( x_1 )</td>
<td>5</td>
<td>17/4</td>
<td>1</td>
<td>0</td>
<td>1/8</td>
<td>-1/8</td>
</tr>
<tr>
<td>( x_2 )</td>
<td>6</td>
<td>19/6</td>
<td>0</td>
<td>1</td>
<td>-1/12</td>
<td>5/12</td>
</tr>
<tr>
<td>( Z_j - C_j )</td>
<td>0</td>
<td>0</td>
<td>1/4</td>
<td>15/8</td>
<td></td>
<td></td>
</tr>
</tbody>
</table>

The solution of the problem \((P_2)\) is \( x_1 = 17/4, x_2 = 19/6, \text{ max } z_2 = 161/4 \)

But \( x_1 \) is an integer value, so for make it integer we use branch and bound method.
Now, the problem \( (P_1) \) is given below:

\[
\begin{align*}
(P_1) & \quad \text{max } z_1 = 5y_1 + 6y_2 \\
& \quad \text{subject to } 10y_1 + 3y_2 \leq 48 \\
& \quad \quad \quad \quad \quad \quad \quad \quad \quad \quad \quad 2y_1 + 3y_2 \leq 12 \\
& \quad \quad \quad \quad \quad \quad \quad \quad \quad \quad \quad y_1 \leq 4 \text{ and integer} \\
& \quad \quad \quad \quad \quad \quad \quad \quad \quad \quad \quad y_2 \leq 10/3 \\
\end{align*}
\]

Similarly using dual simplex method, we obtain optimal solution of \( (P_1) \)

\[y_1 = 9/2, \ y_2 = 1 \text{ and max } z_1 = 57/2.\]

For making \( y_1 \) to be an integer, we use branch and bound method and obtain the solution

\[y_1 = 4, \ y_2 = 4/3 \text{ and max } z_1 = 28\]

Now, the problem \( (P_3) \) is given below:

\[
\begin{align*}
(P_3) & \quad \text{max } z_3 = 5t_1 + 6t_2 \\
& \quad \text{subject to } 10t_1 + 3t_2 \leq 48 \\
& \quad \quad \quad \quad \quad \quad \quad \quad \quad \quad \quad 2t_1 + 3t_2 \leq 12 \\
& \quad \quad \quad \quad \quad \quad \quad \quad \quad \quad \quad t_1 \leq 4 \text{ and integer } t_2 \leq 4/3 \\
\end{align*}
\]

Similarly using dual simplex method and branch and bound method, we can obtain the optimal solution of \( (P_3) \) is

\[t_1 = 4, \ t_2 = 4/3 \text{ and max } z_3 = 28\]

Therefore, the solution for the given fuzzy mixed integer linear programming problem is

\[
\begin{align*}
\bar{x}_1 &= (y_1, x_1, t_1) \quad = \quad (4,4,4) \\
\bar{x}_2 &= (y_2, x_2, t_2) \quad = \quad (4/3, 10/3, 4/3) \\
\bar{z} &= (z_1, z_2, z_3) \quad = \quad (28, 40, 28)
\end{align*}
\]

6. CONCLUSION:

The decomposition method provides an optimal solution to FMILP problems without using ranking functions and applying classical mixed integer linear programming. This method can serve managers by providing the best solution to a variety of mixed integer linear programming problems with fuzzy variables in a simple and effective manner.

REFERENCES:


************